# Symmetric Key based Encryption and Decryption using Lissajous Curve Equations 

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## 1. INTRODUCTION

Symmetric key makes use of same key for encryption and decryption. Symmtric key is confidential for encrypting a plaintext and decryption of plaintext.This application makes use of lissajous cuve equations for key. Plaintext is combined with key for encryption using $16 \times 16$ square matrix. Applying playfair [4] rules for encryption and inverse rules for decryption with same square matrix. Symmetric key playfair cipher is an alphabetic cipher. This application makes use of alphanumeric and special characters for encryption and decryption.

## 2. RESEARCH METHOD

The methodology used in this application is:
a. Key exchange mechanism.
b. Encryption and decryption

In this application parametric equations are used for secure exchange of keys. Lissajous curves equation can be represented ascan be used for two users need to exchange private keys $a$ and $b$ be two large prime numbers' can take values from trigonometric table values and $\mathrm{t}=1$. This is a prerequisite be followed by a sender and recipient.

## 4. User A KEY generation: $a=35761$

```
b=35761
\mp@subsup{k}{x}{}}=30\mathrm{ (in radians)
t=1
ky=120
```

secret key $=38502543$
$\mathrm{dy} / \mathrm{dx}=\left(-\mathrm{b} * \cos \left(\mathrm{k}_{\mathrm{y}}{ }^{*} \mathrm{t}\right) * \mathrm{k}_{\mathrm{y}}\right) /\left(\mathrm{a}^{*} \sin \left(\mathrm{k}_{\mathrm{x}} * \mathrm{t}\right) * \mathrm{kx}\right) \quad$ substituting above values into the equation.

User B KEY generation：
$\mathrm{a}=34551$
b＝34551
$\mathrm{k}_{\mathrm{x}}=30$（in radians）
$\mathrm{t}=1$
$\mathrm{k}_{\mathrm{y}}=120$
$\mathrm{dy} / \mathrm{dx}=\left(-\mathrm{b} * \cos \left(\mathrm{k}_{\mathrm{y}}{ }^{*} \mathrm{t}\right) * \mathrm{k}_{\mathrm{y}}\right) /\left(\mathrm{a}^{*} \sin \left(\mathrm{k}_{\mathrm{x}} * \mathrm{t}\right) * \mathrm{kx}\right) \quad$ substituting above values into the equation．
$\frac{d y}{d x}=1$

Exchange the values of user A＇s kx value to user B and User B‘s ky value to user A．When substituted the value of kx and ky common secret key is generated．Both users arrive at a common number．This common secret key can be further used for encryption and decryption．

| Key | 3 | 8 | 5 | 0 | 2 | 5 | 4 | 3 | 3 | 8 | 5 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Plain text | e | n | e | m | y | k | i | 1 | 1 | e | d |

Figure 1 shows ASCII Key Matrix

|  |  |  |  | ${ }^{1} 15$ |  |  |  |  |  |  |  |  |  |  | 13114 |  | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  |  |  | 田 | ＋ | ${ }^{+1}$ |  |  | 10 | 0 | ［ | 回 | ${ }^{1} 1$ | 이 P | S 月 |  | 翌 |
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| 32 |  | ！ |  | ＂\＃ | \＃ 5 | \％$\%$ | $\varepsilon$ | \＆ |  | 1 | ）＊ | ＊＋ | ＋ | ，－ | － | ． | 7 |
| 48 | 0 | 1 | 2 | 23 | 34 | 45 | 6 | 67 | 78 | 8 | 91 | ： | ＜ | ＜$=$ | ＞ | $>$ | ？ |
| 64 | ［10］ | A | B | B C | C 1 | D E | F | F 6 | H | H | J | J K | K L | L M | 1 H | N | 0 |
| 80 | P | Q | Q R | R S | S 7 | TU | U $V$ | V W | $\mathrm{H}^{2}$ | X Y | Y | z | ［ | ］ | ］ | $\wedge$ |  |
| 96 |  | a | b | $b-$ | c d | de |  | F 9 | 9 h | h | － | j $k$ | k I | 1 m | m $n$ | n | － |
| 118 | P | 9 | 9 | $r \leq$ | $s t$ | t u | 0 | v | － | $\times$ | $y z$ | z | ＜ 1 | $1)$ | － | $\sim$ | 0 |
| 188 | 9 | ij | ie | ¢ 0 | ¢ ${ }^{\text {a }}$ | äd | d | a | E | S | 并这 | － | İ | i | H | － | － |
| 144 | E | 天 | 任 | ff 6 | 6 6 | \％${ }^{\text {d }}$ | di | $\mathrm{Ci}_{4}$ | ij | i | － | 0 | W 5 | $\pm 9$ | ${ }^{1} \mathrm{P}$ P | $\mathrm{F}_{\mathrm{t}}$ | $f$ |
| 150 | G | 1 | 6 | 6 ú | $\square_{4}$ | กัก |  | 9 9 | $2 \dot{1}$ | － | r－ | 2 | $1 / 2$ | 1 | ＊ | ＊ | $\stackrel{ }{*}$ |
| 176 |  | 淴 | 俎 | ，I | $1-1$ | －$\ddagger$ |  |  |  | 7 | 1 | 11 | 7 1 | －ل | $\pm 1$ | $\pm$ |  |
| 198 |  | $\perp$ | ${ }^{\text {T }}$ | T 1 | ＋ | －＋ |  |  | It 4 | ${ }^{4} \mathrm{r}$ | 「 ${ }^{-1}$ | $\xrightarrow{7}$ | तr | $\mathrm{if}^{\text {a }}$ | ＝ | H | $\pm$ |
| 208 | $\Perp$ | $\bar{\sim}$ | \％$\pi$ | $\pi$ u | u | E F |  | $\pi$ \＃ | \＃$\ddagger$ | キ | J | r | － | －1 | I |  | $\square$ |
| 284 | $\alpha$ | F | $\Gamma$ | $\Gamma \pi$ | $\pi \Sigma$ | $\Sigma \sigma$ |  | ${ }^{-1}$ | $r$ r ${ }^{\text {c }}$ | 至 | $8 \Omega$ | $\Omega$ ת | 51 | $\pm 9$ | $\epsilon$ | ¢ | $\square$ |
| 240 |  |  | $\underline{2}$ | $\geq \leq$ | $\leq 1$ |  |  |  |  |  |  |  | 51 | ， |  |  | ， |

Figure 1．ASCII Key Matrix

Encryption．
Encryption process makes use of ASCII key table or square matrix without repetition and characters are arranged in a 16 X 16 grid．

Symmetric Key Based Encryption and Decryption using Lissajous Curve Equations（Santhosh Kumar B．J．）

## Rules:

1. Plaintext letter pair repeated in row are separated with a filler letter, such as $x$ to $l x, l o$ so on.
2. Pick element to right of each letter and left wrap if required. For example, CD is encrypted as DE.
3. Pick element to below each letter and top wrap if required. For example, mu is encrypted as CM.
4. Otherwise, pick same rows opposite corners. The element forms corners of a rectangle. For example BRbecomes Rb.
Decryption is reverse (opposite) of the last rules.
Encrypting a message: "enemy killed"

| Key | 3 | 8 | 5 | 0 | 2 | 5 | 4 | 3 | 3 | 8 | 5 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Plain text | e | n | e | m | y | k | i | 1 | 1 | e | d |
| Cipher text | 5 c | $>\mathrm{h}$ | Eu | $\mathrm{M}^{\prime}$ | 9 r | $; \mathrm{e}$ | 9 d | Lc | Lc | 5 h | 4 e |

Decryption:
Decryption is reverse of encryption.
Considering the ASCII key matrix for decryption. " 5 c " lies in the corners of the rectangle. The order is important - the first character of the encrypted pair is the one that lies on the same row as the first character of the cipher text pair. So " 5 c " can be decrypted as " 3 e " and so on. You will be getting the key with plain text.

| Cipher text | 5 c | $>\mathrm{h}$ | Eu | M' | 9 r | ;e | 9 d | Lc | Lc | 5 h | 4 e |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Key | 3 | 8 | 5 | 0 | 2 | 5 | 4 | 3 | 3 | 8 | 5 |
| Plain text | e | n | e | m | y | k | i | 1 | 1 | e | d |

## 3. RESULTS AND ANALYSIS

In this section, it is explained the results of research and at the same time is given the comprehensive discussion. Results can be presented in figures, graphs, tables and others that make the reader understand easily [2], [5]. The discussion can be made in several sub-chapters.

## 4. CONCLUSION

In this application, new requirement of symmetric key exchange with lissajous curve equations. This application also makes of exchanging secret key between two users using symmetric scheme. Using secret key generated can be used for encryption and decryption. This application makes use of novel method and high security cryptography technique using lissajous curve equations. To provide integrity of any message, message authentication scheme can be applied. It helps to ensure the information secure. Many researches are yet to be identified in future. To provide a better mechanism using asymmetric key exchange lile RSA or DHKE.

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