3283

Minimizing Energy Consumption by Task Consolidation in Cloud Centers with Optimized Resource Utilization

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ABSTRACT

Cloud computing is an emerging field of computation. As the data centers consume large amount of power, it increases the system overheads as well as the carbon dioxide emission increases drastically. The main aim is to maximize the resource utilization by minimizing the power consumption. However, the greatest usages of resources does not mean that there has been a right use of energy. Various resources which are idle, also consumes a significant amount of energy. So we have to keep minimum resources idle. Current studies have shown that the power consumption due to unused computing resources is nearly 1 to 20%. So, the unused resources have been assigned with some of the tasks to utilize the unused period. In the present paper, it has been suggested that the energy saving with task consolidation which has been saved the energy by minimizing the number of idle resources in a cloud computing environment. It has been achieved far-reaching experiments to quantify the performance of the proposed algorithm. The same has also been compared with the FCFSMaxUtil and Energy aware Task Consolidation (ETC) algorithm. The outcomes have shown that the suggested algorithm surpass the FCFSMaxUtil and ETC algorithm in terms of the CPU utilization and energy consumption.

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1. INTRODUCTION

Data communications plays a crucial element of our daily lives and also in the computation. Most of our practices depends on gathering information through the client-server and distributed paradigm [1]. Now a days, client demands have tremendously increased in terms of the number of requests. To cater to the consistent amount of requests, computational capacities and facilities must be constantly reviewed and improved. To remain competitive, the proportional nonnegligible amount of the required energy has been often left behind.

Recent advocacy and tightly coupled with energy consumption called green or sustainable computing has received considerable attention in the field of computation as there is a huge need of minimizing the energy consumption. The scope of sustainable computing goes beyond the main computing

3284 🗖 ISSN: 2088-8708

components, expanding into a much larger range of resources associated with auxiliary equipment, such as the water used for cooling and the physical/floor space occupied by the resources [14].

In Cloud computing, energy consumption and resource utilization are strongly coupled. Specifically, resources with a low utilization rate still consume an unacceptable amount of energy compared to the energy consumption of a fully utilized or sufficiently loaded Cloud computing [10]. According to recent studies [2-5] average resource utilization in most data centers can be as lowas 20%, and the average energy consumption of idle resources can be as high as 60% (or peak power) [3]. To increase resource utilization, task consolidation is an effective technique, greatly enabled by virtualization technologies, which facilitate the concurrent execution of several tasks and, in turn, reduce energy consumption [4].

Power management has been broadly classified into static and dynamic management. Static power anagement has been dealing with fixed power and dynamic power managements with dynamic behaviors for additional degree of capability in virtualized data centers [6]. Infrastructure-as-a-Service (IaaS), Platform-as-a-Service (PaaS), Software-as-a-Service (SaaS) and Database-as-a-Service (DaaS) are the four levels of access in which clouds are deployed for the clients. The task has originated by the different type of customers according to their requirements. There are several heuristic algorithms proposed for local cloud for the centralized controller which has been power aware. Based on the system structure and the characteristics of the cloud infrastructures, a function between the resources of cloud and the combinatorial allocation task has been proposed, as an economics based optimization model [13].

The virtualization concept encapsulates the numerous services that have met the user needs in cloud computing scenario [7]. VMs have been designed to run on various servers which provide the multiple Operating system environments for different applications. Particularly, executing an application which requires resources has been made available for resource provisioning and VM provisioning. Resource provisioning is scheduling the requests for the physical resources where-as VM provisioning creates the instance of VMs as required by the different applications [8].

Server or workload consolidation is the main aim of the task consolidation problem. It allows the servers on a single physical server for minimization of energy consumed by a cloud data center. In the present paper, the task consolidation problem has been addressed to assign n tasks to a set of different resources and the utilization of nodes and distributed VMs are maintained by energy efficiency and load management. The availability of computer nodes during the power consumed by the cloud is the prime concern of the developed algorithm [9-10].

In this paper, the greedy heuristic algorithm has been evaluated and implemented for three basic task consolidations which assign tasks to the physical servers for minimizing the total energy consumed. The proposed heuristic is to minimize the number of idle VMs and minimize the number of idle VMs to as minimum as possible. It has also been shown that the performance improvement is based on different tasks. Section 2 defines the general model of cloud computing environment, energy consumption and task model of the system. We have firmly defined the problem of energy minimization based on the system model. Section 3 deals with the used heuristic algorithm and the algorithms are illustrated by means of an example in Section 4. Section 5 illustrates the set up for the simulation and analyzes the results generated by simulation. Finally, the conclusions have been described in Section 6.

2. CLOUD SYSTEM MODEL

The current section depicts the cloud and its function with the energy models. It also defines the job consolidation problem. The high level architecture of the cloud system is shown in Figure 1 [11]. Virtualization allows the cloud providers to create a set of VMs on a single physical machine that improves the Return on Investment (ROI). The energy consumption may be reduced by switching off the idle nodes, which eliminates the idle power consumption of the given system [12].

In the present work, the target system has been used which consists of a set N of r resources which can be interconnected in the sense that a common route exists between whichever two individual resources as shown in Figure 2. It assumes that the resources are identical in terms of their potential of computing. The virtualization technologies justifies this. The present study has however not considered the federated cloud environment in which the data centers can be placed at different physical locations and the client requests may be processed at various geographical locations.

Figure 3 shows the energy consumption has risen by 56% by data centers from 2005 to 2010 worldwide. Furthermore, CO₂ emissions of the ICT industries are currently approximated to be 2% of the global emissions. It has been observed that global emissions are equivalent to the emissions of the aviation industries. The energy model is conceptualized on the basis that energy consumption has a linear relationship

with processor utilization [11].

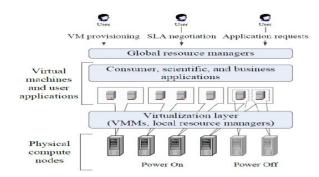


Figure. 1. Cloud System Architecture [11]

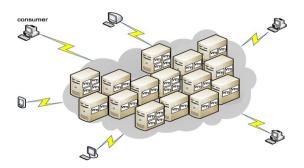


Figure. 2. The cloud Model [8]

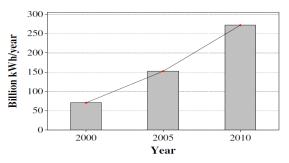


Figure 3. The Worldwide Data Centre Energy Consumption (2000-2010)[11]

This means that for a particular task, the processing time of a task and the processor utilization are the required parameters to determine the energy consumption for that task. The utilization Z_i is defined as for a resource r_i at any given time,

$$Z_i = \sum_{j=1}^t z_{i,j} \tag{1}$$

In (1) it has been observed that t is assigned as the number of tasks running at the current time where as stands for the resource usage of the task $t_{\mathcal{K}_{i,J}}$. The energy consumption E_i of a resource r_i at a given time is derived as

$$E_i = (\delta_{\text{max}} - \delta_{\text{min}}) \times Z_i + \delta_{\text{min}}$$
 (2)

3286 🗖 ISSN: 2088-8708

In (2) it has assigned as δ_{\max} is meaning for the energy consumption at 100% utilization or peak load and δ_{\min} is assigning the minimum energy consumption as low as 1% consumption or in dynamic mode. In this paper, it has assumed that the resources in the objective arrangement are comprised of an efficient power saving method in favor of an inactive time slot. Particularly, the energy use of an inactive source at any specified time has been set to 10% of δ_{\min} . According to the energy consumption, the VMs can be broadly classified into six levels, the idle state and other five levels of CPU utilizations which has shown in Figure 4. In the present study, the task consolidation algorithm assigns a set M of m tasks to a set N of n cloud resources without violating the time constraints to minimize energy consumption and to maximize resource utilization.

```
E_t(V_i) = \begin{cases} \alpha \text{ watts/s, if idle} \\ \beta + \alpha \text{ watts/s, if } 0\% < \text{CPU utilization} \le 50\% \\ 2\beta + \alpha \text{ watts/s, if } 50\% < \text{CPU utilization} \le 70\% \\ 3\beta + \alpha \text{ watts/s, if } 70\% < \text{CPU utilization} \le 80\% \\ 4\beta + \alpha \text{ watts/s, if } 80\% < \text{CPU utilization} \le 90\% \\ 5\beta + \alpha \text{ watts/s, if } 90\% < \text{CPU utilization} \le 100\% \end{cases}
```

Figure 4. Five Levels of CPU Utilization [9]

3. TASK CONSOLIDATION HEURISTIC ALGORITHM

Task allocation is a NP-Hard problem in the cloud. Heuristic and meta-heuristic algorithms are the two useful and efficient technologies for scheduling in cloud due to the ability to distribute and deliver the optimized solutions. In this section, we present the proposed task consolidation algorithm and the performance of the proposed algorithm has been compared with the existing ECTC and FCFSMaxUtil energy conscious task consolidation algorithms.

3.1 FCFSMaxUtil with Minimization of Idle VM

As the idle resources also consume power, the proposed algorithm always minimizes the number of idle resources by allocating a task at an instance to a VM which is currently idle. If no machine is idle it implements FCFSMaxUtil algorithm. The pseudo codes of all the algorithms are as follows:

3.2 ECTC with Minimization of Idle VM

As the idle resources also consume power, the proposed algorithm always minimizes the number of idle resources by allocating a task at an instance to a VM which is currently idle. If no machine is idle it implements ECTC algorithm. The pseudo codes of all the algorithms are as follows:

FCFSMaxUtil Algorithm

```
Input: Task Matrix (TaskId, Arrival Time,
                    Utilization, Execution Time)
Output: Allocation Table (Task Id, Machine Id, Start Time, End Time, Utilization)
[minArrivalTime maxArrivalTime]=
FindMinimumArrivalTimeMaximumArrivalTime(mat)
time= minArrivalTime
while (time <= maxArrivalTime)
do
        Tasklist = GetTasksatArrivalTime(mat,
                  time)
        sort the tasklist in ascending order of arrival
        for each task in tasklist
         do
                 find the VM which has currently
                   highest CPU Utilization
                 Assign the task to the VM and
```

```
update the Allocation table
                 time = time + 1
End Algorithm
ECTC Algorithm
Input: Task Matrix (TaskId, Arrival Time,
                    Utilization, Execution Time)
Output: Allocation Table (Task Id, Machine Id, Start Time, End Time, Utilization)
for each task in tasklist do
        for each vm in vmlist do
                 max=-1
         EnergyConsumedIncludingTheTask(task,vm)
         //Allocate the task into the VM to the
                  maximum energy efficient
        if E > max
                 max=E
                 allocatedVm =Vm
        end if
        end for
        if allocatedVm !=NULL
                 allocate task to allocated VM
                 Update the Alloc table
        end if
end for
End Algorithm
FCFSMaxUtil with Minimization of Idle VM
Input: Task Matrix (TaskId, Arrival Time,
                    Utilization, Execution Time)
Output: Allocation Table (Task Id, Machine Id, Start
  Time, End Time, Utilization)
[minArrivalTime maxArrivalTime]=
FindMinimumArrivalTimeMaximumArrivalTime(mat)
time= minArrivalTime
while (time <= maxArrivalTime)
do
        Tasklist = GetTasksatArrivalTime(mat, time)
        sort the tasklist in ascending order of arrival
         for each task in tasklist do
        for each Vm in Vmlist do
                 if CPUUtilization(Vm)==0
                          allocatedVm=Vm
                          return
                 else
                          find the VM which has
                          currently highest CPU
                          Utilization
                          Assign the task to the
                            VM and update the
                            Allocation table
                 endif
                 time = time + 1
End Algorithm
ECTC with Minimization of Idle VM
Input: Task Matrix (TaskId, Arrival Time,
                    Utilization, Execution Time)
Output: Allocation Table (Task Id, Machine Id, Start
```

```
Time, End Time, Utilization)
for each task in tasklist do
        for each Vm in Vmlist do
                 if CPUUtilization(Vm)==0
                         allocatedVm=Vm
                         return
                 else
                         max = -1
                 E= EnergyConsumedIncludingTheTask(task,Vm)
                 //Allocate the task into the VM to
                 which it will be maximum energy
                 efficient
                 if E > max
                 max=E
                   allocatedVm =Vm
                   end if
end for
if allocatedVm !=NULL
allocate task to allocatedVM
Update the Alloc table
end if
        end for
```

4. AN ILLUSTRATION

Consider a set of 10 VMs $V=\{V1,V2,...V10\}$ and a set of 20 independent tasks $T=\{T1,T2,T3,...T20\}$ in which each task T_i has 4 tuples $\{TaskId, Arrival Time, Processing Time, CPU Utilization\}$. We have considered the threshold value of CPU utilization as 100%. The task table has been shown in Table 1. The task allocation table for various algorithms such as MaxMaxUtil, ECTC and our proposed algorithm are shown in Tables 2, 3,4 and 5 respectively.

Table 1. Example of Task Table

| Table 1. Example of Task Table | | | | | | |
|--------------------------------|---------|------------|-------------|--|--|--|
| Task | Arrival | Processing | Utilization | | | |
| Id | Time | Time | | | | |
| 1 | 1 | 12 | 54 | | | |
| 2 | 1 | 5 | 62 | | | |
| 3 | 1 | 7 | 31 | | | |
| 4 | 1 | 12 | 51 | | | |
| 5 | 1 | 9 | 67 | | | |
| 6 | 2 | 8 | 59 | | | |
| 7 | 2 | 11 | 57 | | | |
| 8 | 2 | 8 | 31 | | | |
| 9 | 2 | 10 | 54 | | | |
| 10 | 2 | 10 | 66 | | | |
| 11 | 2 | 17 | 61 | | | |
| 12 | 3 | 17 | 45 | | | |
| 13 | 3 | 13 | 43 | | | |
| 14 | 3 | 9 | 59 | | | |
| 15 | 3 | 7 | 13 | | | |
| 16 | 3 | 12 | 40 | | | |
| 17 | 4 | 12 | 63 | | | |
| 18 | 4 | 11 | 22 | | | |
| 19 | 4 | 6 | 18 | | | |
| 20 | 4 | 14 | 33 | | | |
| | | | | | | |

Table 2. Task Allocation Table using FCFSMaxutil Algorithm

| T CI SIVILACIII 7 II SOTILIIII | | | | | | |
|--------------------------------|---------|-------|------|-------------|--|--|
| Task | Machine | Start | End | Utilization | | |
| Id | Id | Time | Time | | | |
| 1 | 1 | 1 | 12 | 54 | | |
| 2 | 2 | 1 | 5 | 62 | | |
| 3 | 3 | 1 | 7 | 31 | | |
| 4 | 4 | 1 | 12 | 51 | | |
| 5 | 5 | 1 | 9 | 67 | | |
| 6 | 6 | 2 | 9 | 59 | | |
| 7 | 7 | 2 | 12 | 57 | | |
| 8 | 3 | 2 | 9 | 31 | | |
| 9 | 8 | 2 | 11 | 54 | | |
| 10 | 9 | 2 | 11 | 66 | | |
| 11 | 10 | 2 | 18 | 61 | | |
| 12 | 2 | 6 | 22 | 45 | | |
| 13 | 3 | 10 | 22 | 43 | | |
| 14 | 5 | 10 | 18 | 59 | | |
| 15 | 7 | 3 | 9 | 13 | | |
| 16 | 6 | 10 | 21 | 40 | | |
| 17 | 8 | 12 | 23 | 63 | | |
| 18 | 2 | 6 | 16 | 22 | | |
| 19 | 4 | 4 | 9 | 18 | | |
| 20 | 9 | 12 | 25 | 33 | | |

Table 3. Task Allocation Table using ECTC
Algorithm

Table 4. Task Allocation Table using the Proposed Minimization of Idle VM Algorithm using ECTC

| Algoriumi | | | Minimization of full VIVI Algorithm using ECTC | | | | | | |
|-----------|---------|-------|--|-------------|------|---------|-------|------|-------------|
| Task | Machine | Start | End | Utilization | Task | Machine | Start | End | Utilization |
| Id | Id | Time | Time | | Id | Id | Time | Time | |
| 1 | 1 | 1 | 12 | 54 | 1 | 1 | 1 | 12 | 54 |
| 2 | 2 | 1 | 5 | 62 | 2 | 2 | 1 | 5 | 62 |
| 3 | 1 | 1 | 7 | 31 | 3 | 3 | 1 | 7 | 31 |
| 4 | 3 | 1 | 12 | 51 | 4 | 4 | 1 | 12 | 51 |
| 5 | 4 | 1 | 9 | 67 | 5 | 5 | 1 | 9 | 67 |
| 6 | 5 | 2 | 9 | 59 | 6 | 6 | 2 | 9 | 59 |
| 7 | 6 | 2 | 12 | 57 | 7 | 7 | 2 | 12 | 57 |
| 8 | 3 | 2 | 9 | 31 | 8 | 8 | 2 | 9 | 31 |
| 9 | 7 | 2 | 11 | 54 | 9 | 9 | 2 | 11 | 54 |
| 10 | 8 | 2 | 11 | 66 | 10 | 10 | 2 | 11 | 66 |
| 11 | 9 | 2 | 18 | 61 | 11 | 8 | 2 | 18 | 61 |
| 12 | 7 | 3 | 19 | 45 | 12 | 1 | 3 | 19 | 45 |
| 13 | 6 | 3 | 15 | 43 | 13 | 4 | 3 | 15 | 43 |
| 14 | 10 | 3 | 11 | 59 | 14 | 3 | 3 | 11 | 59 |
| 15 | 1 | 3 | 9 | 13 | 15 | 5 | 3 | 9 | 13 |
| 16 | 10 | 3 | 14 | 40 | 16 | 7 | 3 | 14 | 40 |
| 17 | 2 | 6 | 17 | 63 | 17 | 2 | 6 | 17 | 63 |
| 18 | 9 | 4 | 14 | 22 | 18 | 9 | 4 | 14 | 22 |
| 19 | 3 | 4 | 9 | 18 | 19 | 5 | 4 | 9 | 18 |
| 20 | 8 | 4 | 17 | 33 | 20 | 10 | 4 | 17 | 33 |

Table 5. Task Allocation Table using the Proposed Minimization of Idle VM Algorithm using MAXUTIL

| | of full vivi Algorithm using MAXOTIL | | | | | | |
|------|--------------------------------------|-------|------|-------------|--|--|--|
| Task | Machine | Start | End | Utilization | | | |
| Id | Id | Time | Time | | | | |
| 1 | 1 | 1 | 12 | 54 | | | |
| 2 | 2 | 1 | 5 | 62 | | | |
| 3 | 3 | 1 | 7 | 31 | | | |
| 4 | 4 | 1 | 12 | 51 | | | |
| 5 | 5 | 1 | 9 | 67 | | | |
| 6 | 6 | 2 | 9 | 59 | | | |
| 7 | 7 | 2 | 12 | 57 | | | |
| 8 | 8 | 2 | 9 | 31 | | | |
| 9 | 9 | 2 | 11 | 54 | | | |
| 10 | 10 | 2 | 11 | 66 | | | |
| 11 | 2 | 6 | 22 | 61 | | | |
| 12 | 3 | 8 | 24 | 45 | | | |
| 13 | 5 | 10 | 22 | 43 | | | |
| 14 | 6 | 10 | 18 | 59 | | | |
| 15 | 7 | 3 | 9 | 13 | | | |
| 16 | 8 | 10 | 21 | 40 | | | |
| 17 | 9 | 12 | 23 | 63 | | | |
| 18 | 3 | 8 | 18 | 22 | | | |
| 19 | 4 | 4 | 9 | 18 | | | |
| 20 | 10 | 12 | 25 | 33 | | | |

5. SIMULATION RESULTS

The behavior of three task consolidation heuristic with 1000 tasks has been simulated here. The tasks are bought out for different groups of VMs with the use of incompatible ETC algorithm [6]. Matlab 2012 software has been used for simulation for 1000 tasks. The tasks arrive at the central server queue with a rate of λ having unlimited buffer size. It has taken the arrival interval between the tasks as 1 and the task arrival rate to be 30 in the present studies. The task consolidation algorithms behaviors are demonstrated for 10 and 15 VMs in Figures 5, 6 respectively. The consumption of energy on 15 VMs by varying the task size from 500 to 1500 has been shown in Figure 7.

3290 □ ISSN: 2088-8708

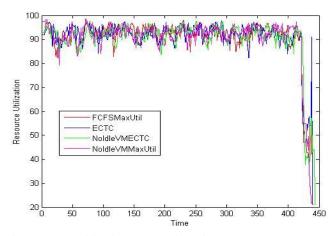


Fig. 5. CPU Utilization Comparison for 1000 Tasks on 10 VMs

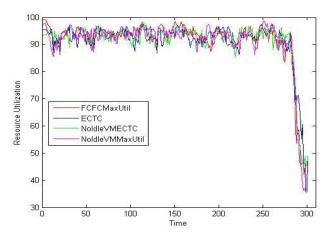


Figure 6. CPU Utilization Comparision for 1000 Taks on 10 VMs

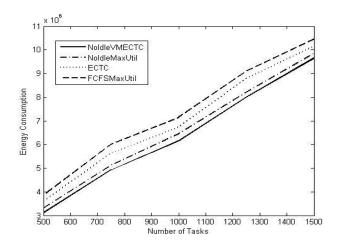


Figure 7. Energy Consumption for Number of Tasks on 15 VMs.

6. CONCLUSION

The simulation experiments have been successfully carried out which examines the behavior of heuristic task consolidation algorithms. It has also been optimized for energy consumption in a cloud

environment. The performance analysis has been demonstrated for the various task consolidation algorithms for the ETC matrix. The results drawn shows that the proposed algorithm has saved the energy as compared to the existing algorithms.

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3292 □ ISSN: 2088-8708



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